

Lecture 19:

Inverse power method with shift

Goal: Take $\mu \in \mathbb{R}$. Find the eigenvalue of A closest to μ .

Observation: Consider $B = A - \mu I$. Then B has eigenvalues:

$$\{\lambda_1 - \mu, \lambda_2 - \mu, \dots, \lambda_n - \mu\} \leftarrow$$

Inverse Power method find eigenvalues such that $|\lambda_j - \mu|$ is the smallest.

$\therefore \lambda_j$ closest to μ can be found.

Algorithm: (Inverse power method with shift)

Step 1: Take $\mu \in \mathbb{R}$. Pick $\vec{x}^{(0)}$ such that $\|\vec{x}^{(0)}\|_\infty = 1$.

Step 2: For $k = 1, 2, \dots$

Solve : $(A - \mu I) \vec{w} = \vec{x}^{(k-1)}$ for \vec{w} .

$$\text{Let : } \vec{x}^{(k)} = \frac{\vec{w}}{\|\vec{w}\|_\infty}.$$

Let $\rho_k = \|A\vec{x}^{(k)}\|_\infty$ ($\rho_k \rightarrow |\lambda_j|$ as $k \rightarrow \infty$)

Recall:

Power's method reads: $\vec{x}^{(k+1)} = \frac{A\vec{x}^{(k)}}{\|A\vec{x}^{(k)}\|_\infty}$ for $k=0, 1, \dots$

$$\Rightarrow \vec{x}^{(k)} = \frac{A^k \vec{x}^{(0)}}{\|A^k \vec{x}^{(0)}\|_\infty}$$

Suppose A is diagonalizable. That's, we can assume $\vec{x}_1, \vec{x}_2, \dots, \vec{x}_n$ form a basis for \mathbb{C}^n .

Take $\vec{x}^{(0)} = a_1\vec{x}_1 + a_2\vec{x}_2 + \dots + a_n\vec{x}_n$ (assuming $a_1 \neq 0$)

$$\vec{x}^{(k)} = \frac{a_1 \lambda_1^k \left[\vec{x}_1 + \sum_{j=2}^n \frac{a_j}{a_1} \left(\frac{\lambda_j}{\lambda_1} \right)^k \vec{x}_j \right]}{\| a_1 \lambda_1^k \left[\vec{x}_1 + \sum_{j=2}^n \frac{a_j}{a_1} \left(\frac{\lambda_j}{\lambda_1} \right)^k \vec{x}_j \right] \|_\infty} \approx \underbrace{\frac{a_1}{\|a_1\|} \frac{\vec{x}_1}{\|\vec{x}_1\|_\infty} \frac{\lambda_1^k}{\|\lambda_1\|^k}}_{\check{v}}$$

(Because $1 > \left| \frac{\lambda_2}{\lambda_1} \right| \geq \left| \frac{\lambda_3}{\lambda_1} \right| \geq \dots \geq \left| \frac{\lambda_n}{\lambda_1} \right|$)


 Rate of convergence of Power's method depends
 on $\left| \frac{\lambda_2}{\lambda_1} \right|$ ← second largest eigenvalue
 ~~$\left| \frac{\lambda_1}{\lambda_1} \right|$~~ ← largest eigenvalue

In fact,

$$\| A \vec{x}^{(k)} \|_\infty \rightarrow \| A \left(\frac{a_1}{\|a_1\|} \frac{\vec{x}_1}{\|\vec{x}_1\|_\infty} \right) \|_\infty = \left\| \frac{a_1}{\|a_1\|} \lambda_1 \frac{\vec{x}_1}{\|\vec{x}_1\|_\infty} \right\|_\infty = \|\lambda_1\|$$

∴

$$\|\lambda_1\| + O\left(\left(\frac{\lambda_2}{\lambda_1}\right)^k\right)$$

Convergence rate: $|\lambda_1| > |\lambda_2| > \dots > |\lambda_n|$

1. Power method:

Converges if $\eta = \left| \frac{\lambda_2}{\lambda_1} \right| < 1$ and $\langle \vec{v}_1, \vec{x}^{(0)} \rangle \neq 0$ (\vec{v}_1 = eigenvector of λ_1)

Also, $\rho_k = \|A\vec{x}^{(k)}\|_\infty = |\lambda_1| + O(\eta^k)$ (Slow convergence if $\eta \approx 1$)

2. Inverse Power method:

Converges if $\left| \frac{1/\lambda_{n-1}}{1/\lambda_n} \right| = \left| \frac{\lambda_n}{\lambda_{n-1}} \right| < 1$ and $\langle \vec{v}_n, \vec{x}^{(0)} \rangle \neq 0$ (\vec{v}_n = eigenvector of λ_n)

Also, $\rho_k = \|A\vec{x}^{(k)}\|_\infty = |\lambda_n| + O(\eta^k)$ (Slow convergence if $\eta \approx 1$)

3. Inverse Power method with shift, let λ_j be closest to μ .

Converges if: $\eta = \max_{m \neq j} \left| \frac{\lambda_j - \mu}{\lambda_m - \mu} \right| < 1$ and $\langle \vec{v}_j, \vec{x}^{(0)} \rangle \neq 0$ (\vec{v}_j = eigenvector of λ_j)

$\rho_k = \|A\vec{x}^{(k)}\|_\infty = |\lambda_j| + O(\eta^k)$ (Slow convergence if $\eta \approx 1$)

How to speed up convergence? Let $A \in M_{n \times n}(\mathbb{R})$

Idea: Use Inverse Power method with shift, update μ in each iteration (such that μ is closer to a real eigenvalue in each iteration)

Then: $\eta \stackrel{\text{def}}{=} \max_{m \neq j} \left| \frac{\lambda_j - \mu}{\lambda_m - \mu} \right|$ becomes smaller and smaller \Rightarrow Converges faster and faster!

Definition: (Rayleigh quotient) Let $\vec{v} \neq \vec{0} \in \mathbb{R}^n$, $A \in M_{n \times n}$. Then, the Rayleigh quotient is defined as: $R(\vec{v}, A) = \frac{\vec{v}^* A \vec{v}}{\vec{v}^* \vec{v}}$.

Remark: Let A be symmetric positive definite. Then: all eigenvalues:

$\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n$ are real.

Then: $\lambda_n \leq R(\vec{v}, A) \leq \lambda_1$ and

$R(\vec{v}, A) = \lambda_1$ when $\vec{v} = \vec{v}_1$ = eigenvector of λ_1 .

$R(\vec{v}, A) = \lambda_n$ when $\vec{v} = \vec{v}_n$ = eigenvector of λ_n

$R(\vec{v}, A)$ can be regarded as the approximation of eigenvalue λ_j , given that \vec{v} is closed to \vec{v}_j .

Rayleigh Quotient Iteration

Let $A \in M_{n \times n}(\mathbb{R})$

Initiate $\vec{x}^{(0)}$ such that $\vec{x}^{(0)} \vec{x}^{(0)*} = 1$

Initiate μ_0 = initial guess of desired eigenvalue.

Solve: $(A - \mu_0 I) \vec{z}_1 = \vec{x}^{(0)}$

Let $\vec{x}^{(1)} = \frac{\vec{z}_1}{\|\vec{z}_1\|_2}$ ($\|\vec{x}\|_2 \stackrel{\text{def}}{=} \sqrt{\vec{x}^* \vec{x}}$)

Let $\mu_1 = R(\vec{x}^{(1)}, A) = \vec{x}^{(1)*} A \vec{x}^{(1)}$ (Improve μ_0 such that it is closer to an actual eigenvalue)
keep iteration going!

Algorithm: (Rayleigh Quotient Iteration)

Input: $\vec{x}^{(0)}$ s.t. $\|\vec{x}^{(0)}\|_2 = 1$ and μ_0 .

Output: μ_k = eigenvalue

For $k=0, 1, 2, \dots$

Step 1: Solve $(A - \mu_k I) \vec{z}_{k+1} = \vec{x}^{(k)}$

Step 2: Let $\vec{x}^{(k+1)} = \frac{\vec{z}_{k+1}}{\|\vec{z}_{k+1}\|_2}$. Step 3: $R(\vec{x}^{(k+1)}, A)$

Example: Let $A = \begin{pmatrix} 1 & 2 & 3 \\ 1 & 2 & 1 \\ 3 & 2 & 1 \end{pmatrix}$.

Eigenvalues: $\lambda_1 = 3 + \sqrt{5}$, $\lambda_2 = 3 - \sqrt{5}$, $\lambda_3 = -2$.

Want to find $3 + \sqrt{5}$.

Let $\vec{x}^{(0)} = \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}$, $\mu_0 = 2.00$

Then: $\vec{x}^{(1)} \approx \begin{pmatrix} -0.57927 \\ -0.57348 \\ -0.57927 \end{pmatrix}$ with $\mu_1 = 5.3355$

Converges very fast!

$$\mu_3 = 5.281 \approx 3 + \sqrt{5}!$$

Remark:

- RQI works for SPP A
- May or may not work for other A.

Lecture 19:

What if λ_1 has multiplicity > 1 ?

Consider the case when A is diagonalizable.

Let $\{\vec{x}_1, \vec{x}_2, \dots, \vec{x}_n\}$ be the basis of eigenvectors with eigenvalues equal to $\lambda_1, \lambda_2, \dots, \lambda_n$.

Assume that: $\lambda_1 = \lambda_2 = \dots = \lambda_i > |\lambda_{i+1}| \geq \dots \geq |\lambda_n|$.

Let $\vec{x}^{(0)} = a_1 \vec{x}_1 + a_2 \vec{x}_2 + \dots + a_n \vec{x}_n$ (with $a_i \neq 0$)

Easy to check: $\vec{x}^{(k)} = \frac{\lambda_1^k (a_1 \vec{x}_1 + \dots + a_i \vec{x}_i + (\frac{\lambda_{i+1}}{\lambda_1})^k \vec{x}_{i+1} + \dots + (\frac{\lambda_n}{\lambda_1})^k \vec{x}_n)}{\|\lambda_1^k (a_1 \vec{x}_1 + \dots + a_i \vec{x}_i + (\frac{\lambda_{i+1}}{\lambda_1})^k \vec{x}_{i+1} + \dots + (\frac{\lambda_n}{\lambda_1})^k \vec{x}_n)\|_\infty}$

$$\rightarrow \frac{a_1 \vec{x}_1 + \dots + a_i \vec{x}_i}{\|a_1 \vec{x}_1 + \dots + a_i \vec{x}_i\|_\infty} \text{ as } k \rightarrow \infty.$$

\uparrow
Eigenvector of λ_1 .

Also, $\|A\vec{x}^{(k)}\|_\infty \rightarrow \left\| \frac{A(a_1 \vec{x}_1 + \dots + a_i \vec{x}_i)}{\|a_1 \vec{x}_1 + \dots + a_i \vec{x}_i\|_\infty} \right\|_\infty = |\lambda_1|$ as $k \rightarrow \infty$

Remark: The condition on multiplicity ($= 1$) can be relaxed.

Method 2: QR method

Preliminary: QR factorization

Definition: $Q \in M_{n \times n}(\mathbb{R})$ is orthogonal if $Q^T Q = I_n$

Remark: - $Q^{-1} = Q^T$

- Columns of Q forms orthonormal set.

Revisit: Gram-Schmidt orthogonalization

Let $A = (\overset{\downarrow}{\vec{a}_1}, \overset{\downarrow}{\vec{a}_2}, \dots, \overset{\downarrow}{\vec{a}_n})$ are linearly independent.

G-S process converts $\{\vec{a}_1, \dots, \vec{a}_n\}$ to orthonormal set $\{\vec{q}_1, \vec{q}_2, \dots, \vec{q}_n\}$

G-S process:

Step 1: Let $\tilde{\vec{q}}_1 = \vec{a}_1$. Normalize : $\vec{q}_1 = \frac{\tilde{\vec{q}}_1}{\|\tilde{\vec{q}}_1\|_2}$. Let $\alpha_{11} = \|\tilde{\vec{q}}_1\|_2$.

Step 2: Define : $\tilde{\vec{q}}_2 = \vec{a}_2 - \alpha_{12} \vec{q}_1$. Choose α_{12} such that :

$$\tilde{\vec{q}}_2^T \tilde{\vec{q}}_1 = 0. \text{ Then: } \alpha_{12} = \vec{q}_1^T \vec{a}_2.$$

Normalize : $\vec{q}_2 = \frac{\tilde{\vec{q}}_2}{\|\tilde{\vec{q}}_2\|_2}$. Let $\alpha_{22} = \|\tilde{\vec{q}}_2\|_2$.

Step 3: Suppose $\vec{q}_1, \vec{q}_2, \dots, \vec{q}_{k-1}$ are constructed.

Let $\tilde{\vec{q}}_k = \vec{a}_k - (\alpha_{1k} \vec{q}_1 + \alpha_{2k} \vec{q}_2 + \dots + \alpha_{(k-1)k} \vec{q}_{k-1})$ where $\alpha_{jk} = \vec{q}_j^T \vec{a}_k$.

Then: $\tilde{\vec{q}}_k^T \vec{q}_i = 0$ for $i = 1, 2, \dots, k-1$. Let $\alpha_{kk} = \|\tilde{\vec{q}}_k\|_2$.

Normalize: $\vec{q}_k = \frac{\tilde{\vec{q}}_k}{\|\tilde{\vec{q}}_k\|_2}$.

In summary,

$$\left\{ \begin{array}{l} \vec{a}_1 = \alpha_{11} \vec{q}_1 \\ \vec{a}_2 = \alpha_{12} \vec{q}_1 + \alpha_{22} \vec{q}_2 \\ \vec{a}_3 = \alpha_{13} \vec{q}_1 + \alpha_{23} \vec{q}_2 + \alpha_{33} \vec{q}_3 \\ \vdots \\ \vec{a}_k = \alpha_{1k} \vec{q}_1 + \alpha_{2k} \vec{q}_2 + \dots + \alpha_{kk} \vec{q}_k \end{array} \right.$$

This is equivalent to:

$$A = \left(\frac{1}{\vec{a}_1}, \frac{1}{\vec{a}_2}, \dots, \frac{1}{\vec{a}_n} \right) = \underbrace{\left(\frac{1}{\vec{q}_1}, \dots, \frac{1}{\vec{q}_n} \right)}_Q \underbrace{\begin{pmatrix} \alpha_{11} & \alpha_{12} & \dots & \alpha_{1n} \\ \alpha_{21} & \alpha_{22} & \dots & \alpha_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ \alpha_{n1} & \alpha_{n2} & \dots & \alpha_{nn} \end{pmatrix}}_R$$

Remark: • Q = orthonormal; R = upper triangular

• Q may NOT be a square matrix ($Q \in M_{m \times n}$)
 R is a square matrix ($R \in M_{n \times n}$)

• Factorization of $A = QR$ is called the QR factorization.